

Session 4 Splitting of fibered categories / Groupoids

§1. Splitting of fibered categories:

Let $F \xrightarrow{\phi} C$ be a fibered category.

• $\forall U \in C$, $F(U)$ a category with obj: $\xi \in F \text{ st } \phi(\xi) = U$

$$\text{maps: } \begin{array}{ccc} \xi & \rightarrow & \delta \\ \downarrow & & \downarrow \\ U & \rightarrow & U \\ & & \text{id}_U \end{array}$$

• $\forall V \rightarrow U$ in C

$\begin{array}{ccc} \uparrow & & \uparrow \\ F(V) \ni \delta & \rightarrow & \xi \in F(U) \end{array}$ st $\delta \rightarrow \xi$ is cartesian.

Want to think of F as a functor

$$\begin{array}{ccc} C^{\text{op}} & \rightarrow & \text{Cat} \\ U & \mapsto & F(U) \end{array}$$

This is not a functor!

If $V \rightarrow U$ in C , do we get a map $F(U) \rightarrow F(V)$?

$$\begin{array}{ccc} V & \xrightarrow{f} & U \\ \uparrow & & \uparrow \\ & & \xi \in F(U) \end{array}$$

Since $F \rightarrow C$ is a fibered category, $\exists \delta \rightarrow \xi$ cartesian lying above f $\xi \in F(U)$

$$\begin{array}{ccc} V & \xrightarrow{f} & U \\ \uparrow & & \uparrow \\ \delta & \rightarrow & \xi \end{array}$$

\leadsto Seeing F as a functor requires choosing cartesian maps lying above maps in C .

Definition (Cleavage):

A cleavage for a fibered category is a choice of a Cartesian map for each map in C and each object in the target.

- Given a fibered category and a cleavage, a map

$f: V \rightarrow U$ induces a functor:

$$f^* := F(f) : F(U) \rightarrow F(V)$$

$$\xi \mapsto \begin{array}{ccc} V & \xrightarrow{f} & U \\ \uparrow & & \uparrow \\ f^*\xi & \longrightarrow & \xi \end{array}$$

$$\delta \mapsto \xi \mapsto f^*\delta \rightarrow \xi$$

- Is $C^{op} \rightarrow \text{Cat}$ a functor?
 $U \mapsto F(U)$
 $f \mapsto f^*$

$$\begin{array}{ccccc} W & \xrightarrow{g} & V & \xrightarrow{f} & U \\ \uparrow & & \uparrow & & \uparrow \\ g^*f^*\xi & \longrightarrow & f^*\xi & \longrightarrow & \xi \\ \in F(W) & & & & \end{array} \qquad \begin{array}{ccc} W & \xrightarrow{f \circ g} & U \\ \uparrow & & \uparrow \\ (f \circ g)^*\xi & \longrightarrow & \xi \\ \in F(W) & & \end{array}$$

$$g^*f^* \cong (f \circ g)^* \\ \text{but } g^*f^* \neq (f \circ g)^* \text{ in general}$$

Definition (Splitting Category)

Let $F \rightarrow C$ be a fibered category. A splitting of F is

a subcategory $K \subset F$ st

1) every arrow in K is Cartesian.

2) $\forall f: V \rightarrow U$ in C and $\forall \xi \in F(U)$, $\exists!$ arrow $\delta \rightarrow \xi$ in K over f .

3) if $U \in \mathcal{C}$ and $\xi \in F(U)$, then $\xi \xrightarrow{id_\xi} \xi$ in $F(U)$ is in \mathcal{K} .

We write (F, \mathcal{K}) a split.

Lemma

The category of split fibered categories that respects the splitting is equivalent to the category $\text{Fun}(\mathcal{C}^{\text{op}}, \text{Cat})$.

Proof:

(\Rightarrow) $\begin{array}{ccc} \mathcal{C}^{\text{op}} & \longrightarrow & \text{Cat} \\ U & \longmapsto & F(U) \\ f & \longmapsto & f^* \end{array}$ is a functor because:

$$\begin{array}{ccc} U & \xrightarrow{id_U} & U & & W & \xrightarrow{g} & V & \xrightarrow{f} & U & \text{ in } \mathcal{C} \\ \uparrow & & \uparrow & & \uparrow & & & & \uparrow & \\ \xi & \xrightarrow{id_\xi} & \xi & & (f \circ g)^* \xi & \longrightarrow & \xi & & & \\ & & & & \parallel & & & & & \\ & & & & g^* f^* \xi & & & & & \end{array}$$

(\Leftarrow) $F: \mathcal{C}^{\text{op}} \rightarrow \text{Cat}$

Let F^{fib} be the fibered category

obj: (U, δ) with $U \in \mathcal{C}$, $\delta \in F(U)$
 maps: $(V, \delta) \rightarrow (U, \delta)$ is a pair (g, α)

$$g: V \rightarrow U \text{ in } \mathcal{C}$$

$$\delta \xrightarrow{\alpha} F(g)(\delta) \text{ in } F(V)$$

$$\begin{array}{ccc} F^{\text{fib}} & \longrightarrow & \mathcal{C} \\ (U, \delta) & \longmapsto & U \end{array}$$

Let K be the subcategory of arrows of the form (g, id)

with g an arrow in C .

$$(*) \quad \begin{array}{ccc} & V & \xrightarrow{f} & U & \text{in } C \\ & \nearrow & & \uparrow & \\ (V, \delta) & \longrightarrow & (U, \delta) & \in F^{\text{fib}}(U) = K(U) \end{array}$$

with $\delta \rightarrow F(f)(\delta)$

$$\Rightarrow \delta = F(f)(\delta) \xrightarrow{\text{id}} F(f)(\delta) \quad \text{for a lift in } K$$

$\Rightarrow (f, \text{id})$ unique possible choice in K .

(*) K contains all identities:

$$U \in C \text{ and } \delta \in F(U)$$

$$(U, \delta) \xrightarrow{(\text{id}, \text{id})} (U, \delta) \text{ is in } K$$

$$\delta \rightarrow F(\text{id})(\delta) = \text{id}(\delta) = \delta \quad \text{because } F \text{ is a functor}$$

(*) All arrows are Cartesian:

$$\begin{array}{ccccc} W & & & & \\ \uparrow & \searrow & \xrightarrow{f} & & \\ h & \rightarrow & V & \xrightarrow{g} & U & \text{in } C \\ & & \uparrow & & \uparrow \\ & & (V, \delta) & \xrightarrow{(g, \text{id})} & (U, \delta) \\ & \nearrow & & & \nearrow \\ (W, \alpha) & \xrightarrow{(f, \text{id})} & & & \end{array}$$

(h, id) (dotted arrow)
 $\delta = F(g)(\delta)$ (red arrow from (U, δ))

$$\begin{aligned} \alpha &= F(f)(\delta) \\ &= F(g \circ h)(\delta) \\ &= F(h) \circ F(g)(\delta) \quad \text{F is a functor} \\ &= F(h)(\delta) \end{aligned}$$

Examples:

1) $F: C^{op} \rightarrow Set$

F^{fib} is a splitting fibered category.

2) G a group seen as a category with one object.

obj: $*_G$
 maps: $*_G \rightarrow *_G$ (elems of G)
 $g \in G$

Let G and H be groups. A functor $f: G \rightarrow H$ is the same as gp homomorphism.

• $G \xrightarrow{f} H$ is fibered category iff f is surjective

$$\begin{array}{ccc}
 *_G & \xrightarrow{g \in G} & *_G & \text{in } G \\
 \downarrow & & \downarrow & \downarrow f \\
 *_H & \xrightarrow{h \in H} & *_H & \text{in } H
 \end{array}$$

st $f(g) = h$

• All arrows are cartesian

$$\begin{array}{ccccc}
 & & & & \xrightarrow{g_1 \in G} \\
 & & & & *_G \\
 & \xrightarrow{g_2 \in G} & & & \downarrow \\
 *_G & \xrightarrow{g_2 g_1^{-1}} & *_G & \xrightarrow{g_1 \in G} & *_G \\
 & & \downarrow & & \downarrow \\
 & & *_H & \xrightarrow{h_1 \in H} & *_H \\
 & \xrightarrow{h_2 = h_1 h_1^{-1} \in H} & & & \downarrow \\
 *_H & & & & *_H \\
 & & & & \xrightarrow{h_2 \in H}
 \end{array}$$

• What is a splitting here?

Assume $f: G \rightarrow H$ surjective so that $G \rightarrow H$ is a fibered category.

a splitting = a section $H \xrightarrow{s} G$ which is gp homomorphism.

To get a non-example, consider any group extension G of H by K that doesn't split.

Theorem:

Let $F \rightarrow C$ be a fibered category.

There exists a split fibered category (\hat{F}, k) st \hat{F} is equivalent to F .

Proof:

Let $H: C^{op} \rightarrow \text{Cat}$

$$H \stackrel{\text{fib}}{=} U \mapsto \text{HOM}_C(C/U, F)$$

Let \hat{F} be the split fibered category associated to the functor H .

$$\begin{aligned} \hat{F}: \text{obj: } (U, \delta) \text{ with } U \in C \text{ and } \delta \in H(U) = \text{HOM}(C/U, F) \\ \text{maps: } (V, \delta) \xrightarrow{(g, \alpha)} (U, \delta) \quad g: V \rightarrow U \text{ in } C \\ \alpha \text{ base preserving nat. transf} \end{aligned}$$

Show that \hat{F} and F are equivalent.

$$\begin{aligned} \hat{F} &\xrightarrow{e} F \\ (U, \delta) &\mapsto \delta(\text{id}_U) \\ (g, \alpha) &\mapsto \delta(g) \circ \alpha(\text{id}_V) \end{aligned}$$

Check that:

- it is a morphism of fibered categories
 - e restricts to the evaluation map on the fibers
- \leadsto 2-Yoneda lemma, e is an equivalence on each fiber $\Rightarrow e$ an equivalence



Remark :

$\text{Fun}(C^{\text{op}}, \text{Cat}) \hookrightarrow$ Category of fibered categories over C
is faithful + essentially surjective but not full.

§ 2. Groupoids / Categories fibered in groupoids:

Definition (Groupoid)

A groupoid is a category where all morphisms are isomorphisms.

Examples:

1) Any set seen as a category.

2) G a group seen as a category with one object.

Definition (Category fibered in groupoids)

- A category fibered in groupoids over a category C is a fibered category $F \rightarrow C$ st $\forall U \in C$, $F(U)$ is a groupoid.

Example:

$G \xrightarrow{f} H$ surjective gp homomorphism.

it is a category fibered in groupoids.

Proposition

let F and F' be fibered categories in groupoids over C .

Then, $\text{HOM}_C(F, F')$ is a groupoid.

Proof:

Let $f, g : F \rightarrow F'$ morphisms of fibered categories

and $\xi : f \rightarrow g$ a morphism. Is it an isomorphism?

For $u \in F$, $\xi_u : f(u) \rightarrow g(u)$ an arrow in F'

Let $X \in C$ the image of u via $p_F : F \rightarrow C$

$$\begin{array}{ccc} f(u) & \xrightarrow{\xi_u} & g(u) & \text{in } F' \\ \downarrow & & \downarrow & \\ X & \xrightarrow{\text{id}_X} & X & \text{in } C \end{array}$$

ξ_u base-pres. $\left(\right.$

hence ξ_u is an arrow in $F'(X)$ which is a groupoid. \square

Definition (fiber product of groupoids)

$$\begin{array}{ccc} & G_1 & \\ & \downarrow f & \\ G_2 & \xrightarrow{g} & G \end{array}$$

f, g functors

a diagram of groupoids. We define a groupoid $G_1 \times_G G_2$ as follows:

obj: triples (u, y, δ) where $u \in G_1$, $y \in G_2$ and δ is a morphism $f(u) \rightarrow g(y)$ in G .

maps: $(u', y', \delta') \rightarrow (u, y, \delta)$ is a pair of isomorphisms

$$u' \xrightarrow{a} u \quad y' \xrightarrow{b} y \quad \text{st}$$

$$\begin{array}{ccc} f(u') & \xrightarrow{\quad} & g(y') \\ f(a) \downarrow & \hookrightarrow & \downarrow f(b) \\ f(u) & \xrightarrow{\quad} & g(y) \end{array} \quad \text{Commutative.}$$

check it is a groupoid.

- There are functors $p_1 : G_1 \times G_2 \rightarrow G_1$ and a natural isomorphism of functors $\Sigma : f \circ p_1 \rightarrow g \circ p_2$

$$\begin{array}{ccc}
 G_1 \times G_2 & \xrightarrow{p_2} & G_2 \\
 \downarrow p_1 & & \downarrow f \\
 G_1 & \xrightarrow{g} & G
 \end{array}$$

The category $G_1 \times G_2$, together with the functors p_1 and p_2 and Σ have the following universal property:

If \mathcal{H} is another groupoid and

$$\begin{array}{ccc}
 \mathcal{H} & \xrightarrow{\beta} & G_2 \\
 \downarrow \alpha & \dashrightarrow h & \downarrow f \\
 G_1 \times G_2 & \xrightarrow{p_2} & G_2 \\
 \downarrow p_1 & & \downarrow f \\
 G_1 & \xrightarrow{g} & G
 \end{array}$$

and $\delta : f \circ \beta \rightarrow g \circ \alpha$ an isomorphism of functors

Then, there exists $(h : \mathcal{H} \rightarrow G_1 \times G_2, \alpha_1, \alpha_2)$

where h is a functor

$$\left. \begin{array}{l}
 \alpha_1 : \mathcal{H} \rightarrow p_1 \circ h \\
 \alpha_2 : \mathcal{H} \rightarrow p_2 \circ h
 \end{array} \right\} \text{ isom of functors}$$

and

$$\begin{array}{ccc}
 f \circ \alpha & \rightarrow & f \circ p_1 \circ h \\
 \delta \downarrow & & \downarrow \Sigma \circ h \\
 g \circ \beta & \rightarrow & g \circ p_2 \circ h
 \end{array}$$

The data (h, α, β) is unique up to a unique isomorphism.

• Let C be a category and

$$\begin{array}{ccc} & F_1 & \\ & \downarrow^c & \\ F_2 & \xrightarrow{\quad} & F_3 \\ & \downarrow^d & \end{array}$$

a diagram of fibered categories in groupoids over C .

• Consider G a category fibered in groupoids over C , with

$$\left. \begin{array}{l} \alpha: G \rightarrow F_1 \\ \beta: G \rightarrow F_2 \end{array} \right\} \text{morphisms of f.c.}$$

$$c \circ \alpha \xrightarrow{\delta} d \circ \beta \quad \text{an isomorphism of fibered cat } G \rightarrow F_3$$

Giving the data (α, β, δ) is equivalent to giving an object of

$$\frac{\text{HOM}_C(G, F_1) \times \text{HOM}_C(G, F_2)}{\text{HOM}_C(G, F)}$$

Proposition:

\nearrow = fiber product category

There exists a collection of data $(G, \alpha, \beta, \delta)$ as above st for every H category fibered in groupoids over C ,

$$\text{HOM}_C(H, G) \rightarrow \frac{\text{HOM}(H, F_1) \times \text{HOM}(H, F_2)}{\text{HOM}(H, F_3)}$$

$$h \mapsto (\alpha \circ h, \beta \circ h, \delta \circ h)$$

is an isomorphism

(IV)

Let G be a groupoid. We can describe it as follows:

- G_0 set of objects
- for each objects x, y , a set $G(x, y)$ of morphisms from x to y .
- for every object x , a designated element id_x in $G(x, x)$.
- for each triple objects x, y and z , a function

$$\text{comp: } G(y, z) \times G(x, y) \rightarrow G(x, z)$$
$$(g, f) \mapsto g \circ f$$

- for each objects x, y , an inverse function

$$G(x, y) \rightarrow G(y, x) : f \mapsto f^{-1}$$
$$x \rightarrow y \quad y \rightarrow x$$

All satisfying, for any $f: x \rightarrow y$, $g: y \rightarrow z$, $h: z \rightarrow w$:

$$f \circ \text{id}_x = f \quad \text{and} \quad \text{id}_y \circ f = f$$

$$(h \circ g) \circ f = h \circ (g \circ f)$$

$$f \circ f^{-1} = \text{id}_y \quad \text{and} \quad f^{-1} \circ f = \text{id}_x$$

If $f \in G(x, y)$, x is called source of f and y target of f .

We write

$$x = s(f)$$
$$y = t(f)$$

Let C be a category with finite fiber products.

A groupoid object in C consists of :

- A pair (U, R) of objects

- Five ^{obj} morphisms ^{arrows} :

$$\begin{array}{ccccc}
 s, t: & R & \rightarrow & U & , & e: U & \rightarrow & R & , & i: R & \rightarrow & R \\
 \swarrow & & & \searrow & & \downarrow & & & & \downarrow & & \\
 \text{source} & & & \text{target} & & \text{unit} & & & & \text{inverse} & &
 \end{array}$$

$$\begin{array}{ccc}
 m: R \times R & \rightarrow & R \\
 \text{multiplication} & & \text{multiplication} \\
 \text{source } s, U, t & &
 \end{array}$$

satisfying the following :

$$\begin{array}{ccc}
 1) & U \xrightarrow{e} R & R \xleftarrow{p_1} R \times R \xrightarrow{p_2} R \\
 & \begin{array}{ccc} \varepsilon \downarrow & \searrow \text{id}_U & \downarrow t \\ R & \xrightarrow{s} & U \end{array} & \begin{array}{ccc} s \downarrow & \hookrightarrow & \downarrow m \hookrightarrow \\ U & \xleftarrow{s} & R \xrightarrow{t} U \end{array}
 \end{array}$$

$$\begin{array}{ccccc}
 \text{(Unit)} & U \times R & = & R & = & R \times U \\
 & \downarrow \text{id}_U, t & & \parallel & & \downarrow s, U, \text{id} \\
 & \text{exists} & & & & \downarrow \text{id}_U \times e \\
 & R \times R & \xrightarrow{m} & R & \xleftarrow{m} & R \times R \\
 & \text{source } s, U, t & & & & \text{source } s, U, t
 \end{array}$$

$$\begin{array}{ccccc}
 \text{(inverse)} & R & \xrightarrow{i \times \text{id}} & R \times R & \xleftarrow{\text{id} \times i} & R \\
 & \downarrow s & & \downarrow m & & \downarrow t \\
 & U & \xrightarrow{e} & R & \xleftarrow{e} & U
 \end{array}$$

(associativity)

$$\begin{array}{ccc} R \times R \times R & \xrightarrow{m \times \text{id}} & R \times R \\ s, U, t & s, U, t & s, U, t \\ \text{id} \times m \downarrow & \curvearrowright & \downarrow m \\ R \times R & \xrightarrow{m} & R \\ s, U, t & & \end{array}$$

Examples:

1) A groupoid object in the category of sets (Set) is a groupoid in the usual sense.

U = the set of all objects in Set

R = the set of all arrows in Set

$$s(a \rightarrow b) = a \quad t(a \rightarrow b) = b$$

$$e(a) = \text{id}_a \quad i \left(\underset{\in R}{f: a \rightarrow b} \right) = \underset{\in R}{f^{-1}: b \rightarrow a}$$

$$m(f, g) = g \circ f.$$

3) Let S be a scheme and G/S a group scheme acting on an S -scheme X . Define the action groupoid associated to this G -action as:

$$(i) \quad U = X, \quad R = X \times_S G$$

(ii) $s: X \times_S G \rightarrow X$ is the first projection.

$t: X \times_S G \rightarrow X$ is the action map.

(iii) $e: X \rightarrow X \times_S G$ is induced by $S \rightarrow G$ section.

(iv) $i: X \times_S G \rightarrow X \times_S G, (u, g) \mapsto (g \cdot u, g^{-1})$

$$(v) \quad X \times_S G \times_X X \times_S G \xrightarrow{m} X \times_S G$$

$$X \times_S G \times_S G$$

is induced by $G \times_S G \rightarrow G$ the law.

then we get a groupoid object in (Sch/S)

usually denoted $\{X/G\}$.

Discussion: difference between a groupoid object in \mathcal{C} and a group object in \mathcal{C} .